Parallel Programming Systems and Models

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CS 595 Hot Topics in Distributed Systems: Data-Intensive Computing October 13th, 2010

- Moore's Law
 - The number of transistors that can be placed inexpensively on an integrated circuit will double approximately every 18 months.
 - Self-fulfilling prophecy
 - Computer architect goal
 - Software developer assumption

- Impediments to Moore's Law
 - Theoretical Limit
 - What to do with all that die
 - Design complexity
 - How do you meet the exp increase?



- von Neumann model
 - Execute a stream of instructions (machine code)
 - Instructions can specify
 - Arithmetic operations
 - Data addresses
 - Next instruction to execute
 - Complexity
 - Track billions of data locations and millions of instructions
 - Manage with:
 - Modular design
 - High-level programming languages

- Parallelism
 - Continue to increase performance via parallelism.



- From a software point-of-view, need to solve demanding problems
 - Engineering Simulations
 - Scientific Applications
 - Commercial Applications
- Need the performance, resource gains afforded by parallelism

- Engineering Simulations
 - Aerodynamics
 - Engine efficiency



Scientific Applications

 Bioinformatics
 Thermonuclear processes
 Weather modeling





- Commercial Applications
 - Financial transaction processing
 - Data mining
 - Web Indexing



- Unfortunately, greatly increases coding complexity
 - Coordinating concurrent tasks
 - Parallelizing algorithms
 - Lack of standard environments and support

- The challenge
 - Provide the abstractions, programming paradigms, and algorithms needed to effectively design, implement, and maintain applications that exploit the parallelism provided by the underlying hardware in order to solve modern problems.

• Standard sequential architecture



- Use multiple
 - Datapaths
 - Memory units
 - Processing units

SIMD

- Single instruction stream, multiple data



SIMD

- Advantages
 - Performs vector/matrix operations well
 - EX: Intel's MMX chip
- Disadvantages
 - Too dependent on type of computation
 EX: Graphics
 - Performance/resource utilization suffers if computations aren't "embarrasingly parallel".

• MIMD

Multiple instruction stream, multiple data stream



• MIMD

- Advantages
 - Can be built with off-the-shelf components
 - Better suited to irregular data access patterns
- Disadvantages
 - Requires more hardware (!sharing control unit)
 - Store program/OS at each processor
- Ex: Typical commodity SMP machines we see today.

- Task Communication
 - Shared address space
 - Use common memory to exchange data
 - Communication and replication are implicit
 - Message passing
 - Use send()/receive() primitives to exchange data
 - Communication and replication are explicit

- Shared address space
 - Uniform memory access (UMA)
 - Access to a memory location is independent of which processing unit makes the request.
 - Non-uniform memory access (NUMA)
 - Access to a memory location depends on the location of the processing unit relative to the memory accessed.

- Message passing
 - Each processing unit has its own private memory
 - Exchange of messages used to pass data
 - APIs
 - Message Passing Interface (MPI)
 - Parallel Virtual Machine (PVM)

- Algorithm
 - a sequence of finite instructions, often used for calculation and data processing.
- Parallel Algorithm
 - An algorithm that which can be executed a piece at a time on many different processing devices, and then put back together again at the end to get the correct result

- Challenges
 - Identifying work that can be done concurrently.
 - Mapping work to processing units.
 - Distributing the work
 - Managing access to shared data
 - Synchronizing various stages of execution.

- Models
 - A way to structure a parallel algorithm by selecting decomposition and mapping techniques in a manner to minimize interactions.

- Models
 - Data-parallel
 - Task graph
 - Work pool
 - Master-slave
 - Pipeline
 - Hybrid

- Data-parallel
 - Mapping of Work
 - Static
 - Tasks -> Processes
 - Mapping of Data
 - Independent data items assigned to processes (Data Parallelism)

- Data-parallel
 - Computation
 - Tasks process data, synchronize to get new data or exchange results, continue until all data processed
 - Load Balancing
 - Uniform partitioning of data
 - Synchronization
 - Minimal or barrier needed at end of a phase
 - Examples
 - Ray Tracing

 Data-parallel 0 Ρ D D ()Ρ D D \cap Ρ D \mathbf{O} Ρ 0 Ρ

- Task graph
 - Mapping of Work
 - Static
 - Tasks are mapped to nodes in a data dependency task dependency graph (Task parallelism)
 - Mapping of Data
 - Data moves through graph (Source to Sink)

- Task graph
 - Computation
 - Each node processes input from previous node(s) and send output to next node(s) in the graph
 - Load Balancing
 - Assign more processes to a given task
 - Eliminate graph bottlenecks
 - Synchronization
 - Node data exchange
 - Examples
 - Parallel Quicksort, Divide and Conquer approaches
 - Scientific Applications that can be expressed in workflows (e.g. DAGs)

 Task graph Ρ P Ρ Ρ 0 Ρ D 0 D D Р

- Work pool
 - Mapping of Work/Data
 - No desired pre-mapping
 - Any task performed by any process
 - Pull-model oriented
 - Computation
 - Processes work as data becomes available (or requests arrive)

- Work pool
 - Load Balancing
 - Dynamic mapping of tasks to processes
 - Synchronization
 - Adding/removing work from input queue
 - Examples
 - Web Server
 - Bag-of-tasks

Work pool



- Master-slave
 - Modification to Worker Pool Model
 - One or more Master processes generate and assign work to worker processes\
 - Push-model oriented
 - Load Balancing
 - A Master process can better distribute load to worker processes

Pipeline

- Mapping of work
 - Processes are assigned tasks that correspond to stages in the pipeline
 - Static
- Mapping of Data
 - Data processed in FIFO order
 - Stream parallelism

Pipeline

- Computation
 - Data is passed through a succession of processes, each of which will perform some task on it
- Load Balancing
 - Insure all stages of the pipeline are balanced (contain the same amount of work)
- Synchronization
 - Producer/Consumer buffers between stages
- Ex: Processor pipeline, graphics pipeline

• Pipeline



- Message-Passing
- Shared Address Space

- Message-Passing
 - Most widely used for programming parallel computers (clusters of workstations)
 - Key attributes:
 - Partitioned address space
 - Explicit parallelization
 - Process interactions
 - Send and receive data

- Message-Passing
 - Communications
 - Sending and receiving messages
 - Primitives
 - send(buff, size, destination)
 - receive(buff, size, source)
 - Blocking vs non-blocking
 - Buffered vs non-buffered
 - Message Passing Interface (MPI)
 - Popular message passing library
 - ~125 functions

Message-Passing



- Shared Address Space
 - Mostly used for programming SMP machines (multicore chips)
 - Key attributes
 - Shared address space
 - Threads
 - Shmget/shmat UNIX operations
 - Implicit parallelization
 - Process/Thread communication
 - Memory reads/stores

- Shared Address Space
 - Communication
 - Read/write memory
 - EX: x++;
 - Posix Thread API
 - Popular thread API
 - Operations
 - Creation/deletion of threads
 - Synchronization (mutexes, semaphores)
 - Thread management

Shared Address Space



Parallel Programming Pitfalls

- Synchronization
 - Deadlock
 - Livelock
 - Fairness
- Efficiency
 - Maximize parallelism
- Reliability
 - Correctness
 - Debugging

Questions

