Problem 1 Let $G = (V, E)$ be a simple undirected graph with weights $w : E \to \mathbb{Z}^+$. The inductivity of a vertex ordering (permutation $\Pi$ of $V$) $\langle v_{\Pi(1)}, v_{\Pi(2)}, \ldots, v_{\Pi(n)} \rangle$ is defined by

$$\max_{2 \leq j \leq n} \sum_{1 \leq i < j} w(v_{\Pi(i)}v_{\Pi(j)}).$$

(1)

Use (even if not completely covered yet) Fibonacci heaps to obtain a $O(|E| + |V| \log |V|)$-time algorithm (present pseudocode) to produce a least-inductivity vertex ordering of $G$, together with the proof of correctness. That is, find the permutation $\Pi$ of $V$ that minimizes Formula (1)

**Hint:** Use a greedy strategy paying attention to nodes with smallest weighted degree in $G$. Give a $O(|E| + |V|)$-time algorithm for the unweighted case (all the weights are 1).

Problem 2 Describe a binary search tree on $n$ nodes such that the average depth of a node in the tree is $\Theta(\lg n)$ but the height of the tree is not $O(\lg n)$. How large can the height of an $n$-node binary search tree be if the average depth of a node is $\Theta(\lg n)$?

Problem 3 Assume every node in a binary search tree has a pointer to its parent, in addition to pointers to the left and right child. Design an algorithm (write pseudocode), which, given a node $v$, finds $w$, the node-successor of $v$ in inorder (the element of $w$ is also the successor of the element of $v$ in the sorted order of elements).

Analyze the running time of $s$ consecutive calls to successor (that is, $w$ is given as the argument to the next call, and so on) in terms of $s$ and $h$, the height of the tree. A tight (within a constant) analysis is worth one third of the points.

Problem 4 Suppose we wish not only to increment a binary number, but also to reset it to zero (i.e., make all bits in it 0). Counting the cost to examine or modify a bit as 1, show how to implement a binary number as an array of bits so that any sequence of $n$ INCREMENT and RESET operations costs $O(n)$ on an initially zero number.

**Hint:** Keep a pointer to the high-order 1.