MapReduce

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Logistics

- Quiz #3 answers
- Project Proposals
- Feedback before Wednesday
 - Some groups might get a 2nd chance to redo proposals
- Will post reading assignments today

Logistics

- Plan for rest of the semester
 - MapReduce/Hadoop
 - Swift
 - Spark/Sparrow/Mesos
 - Ceph
 - Lustre/GPFS/PVFS
 - More to come

Motivation: Large Scale Data Processing

Want to:

- Process lots of data (> 1 TB)
- Automatically parallelize across hundreds/ thousands of CPUs
- Have status and monitoring tools
- Provide clean abstraction for programmers
- Make this easy

MapReduce

 "A simple and powerful interface that enables automatic parallelization and distribution of large-scale computations, combined with an implementation of this interface that achieves high performance on large clusters of commodity PCs."

Typical Problem

- Map Iterate over a large number of records
 - Extract something of interest from each
 - Shuffle and sort intermediate results
 - Aggregate intermediate results reduce
 - Generate final output

 Key idea: provide an abstraction at the point of these two operations

MapReduce: Programming Model

- Process data using special map() and reduce() functions
- The map() function is called on every item in the input and emits a series of intermediate key/ value pairs
- All values associated with a given key are grouped together
- The reduce() function is called on every unique key, and its value list, and emits a value that is added to the output

Programming Model

- Borrows from functional programming
- Users implement interface of two functions:

```
- map (in_key, in_value) ->
    (out_key, intermediate_value) list
- reduce (out_key, intermediate_value list) ->
    out_value list
```

map

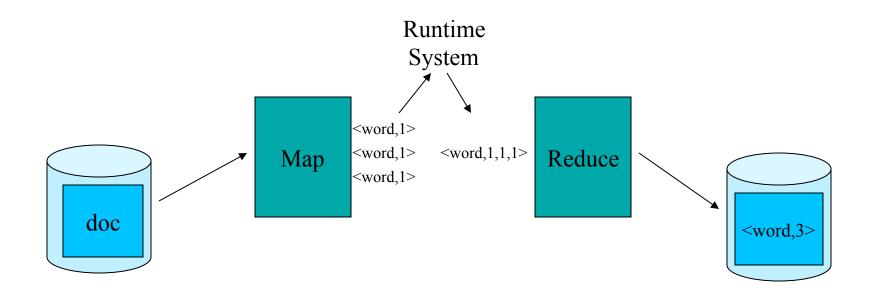
- Records from the data source (lines out of files, rows of a database, etc) are fed into the map function as key*value pairs: e.g., (filename, line).
- map() produces one or more intermediate values along with an output key from the input.

reduce

- After the map phase is over, all the intermediate values for a given output key are combined together into a list
- reduce() combines those intermediate values into one or more final values for that same output key
- (in practice, usually only one final value per key)

MapReduce Examples

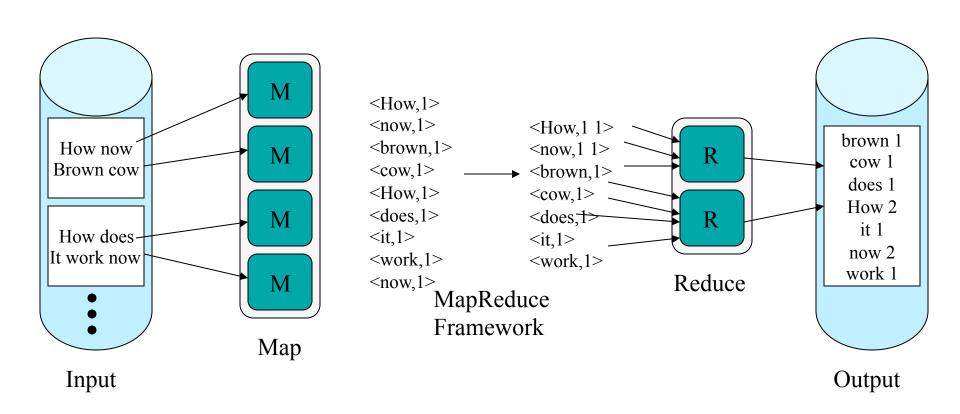
Word frequency



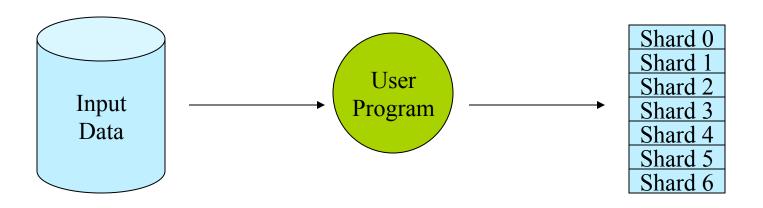
MapReduce Examples

- Distributed grep
 - Map function emits <word, line_number> if word matches search criteria
 - Reduce function is the identity function
- URL access frequency
 - Map function processes web logs, emits <url, 1>
 - Reduce function sums values and emits <url, total>

MapReduce: Programming Model

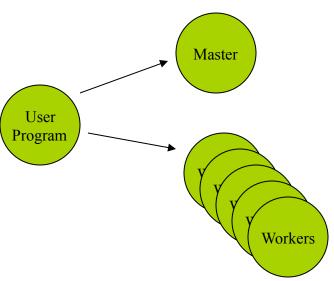


1. The user program, via the MapReduce library, shards the input data



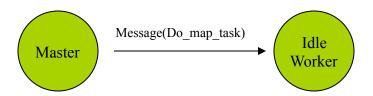
^{*} Shards are typically 16-64mb in size

 The user program creates process copies distributed on a machine cluster.
 One copy will be the "Master" and the others will be worker threads.



MapReduce Resources

- 3. The master distributes M map and R reduce tasks to idle workers.
 - M == number of shards
 - R == the intermediate key space is divided into R parts

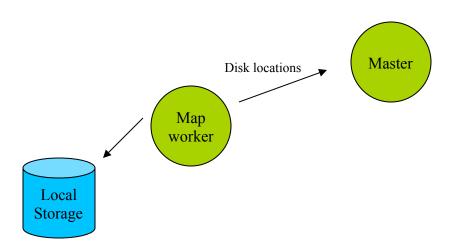


MapReduce Resources

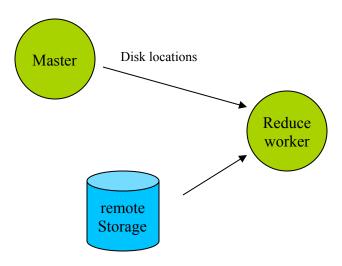
- Each map-task worker reads assigned input shard and outputs intermediate key/value pairs.
 - Output buffered in RAM.



 Each worker flushes intermediate values, partitioned into R regions, to disk and notifies the Master process.



 Master process gives disk locations to an available reduce-task worker who reads all associated intermediate data.

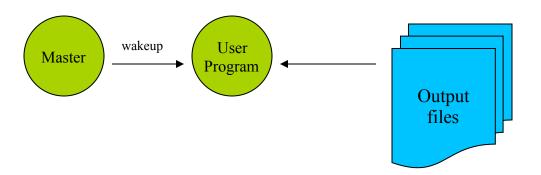


7. Each reduce-task worker sorts its intermediate data. Calls the reduce function, passing in unique keys and associated key values. Reduce function output appended to reduce-task's partition output file.

Sorts data

Partition Output file

8. Master process wakes up user process when all tasks have completed. Output contained in R output files.



MapReduce Runtime System

- 1. Partitions input data
- 2. Schedules execution across a set of machines
- 3. Handles machine failure
- 4. Manages interprocess communication

Parallelism

- map() functions run in parallel, creating different intermediate values from different input data sets
- reduce() functions also run in parallel, each working on a different output key
- All values are processed independently
- Bottleneck: reduce phase can't start until map phase is completely finished.

Locality

- Master program divvies up tasks based on location of data: tries to have map() tasks on same machine as physical file data, or at least same rack
- map() task inputs are divided into 64 MB blocks: same size as Google File System chunks

Fault Tolerance

- Master detects worker failures
 - Re-executes completed & in-progress map() tasks
 - Re-executes in-progress reduce() tasks
- Master notices particular input key/values cause crashes in map(), and skips those values on re-execution.
 - Effect: Can work around bugs in third-party libraries!

Optimizations

- No reduce can start until map is complete:
 - A single slow disk controller can rate-limit the whole process
- Master redundantly executes "slowmoving" map tasks; uses results of first copy to finish

MapReduce Conclusions

- MapReduce has proven to be a useful abstraction
- Greatly simplifies large-scale computations at Google
- Functional programming paradigm can be applied to large-scale applications
- Fun to use: focus on problem, let library deal w/ messy details
- Greatly reduces parallel programming complexity
 - Reduces synchronization complexity
 - Automatically partitions data
 - Provides failure transparency
 - Handles load balancing

Hadoop

- Open source MapReduce implementation
 - http://hadoop.apache.org/core/index.html
- Uses
 - Hadoop Distributed Filesytem (HDFS)
 - http://hadoop.apache.org/core/docs/current/ hdfs design.html
 - Java
 - -ssh

Questions

