Large Scale Wireless Network Systems: Experience, Observations, and Theories

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Wireless Sensor/Actuator Networks

Bridging the digital world and physical world



Wide Applications: CPS, IOT



Environment



Transportation



Smart Grid



Security



Green Building



Industry Monitoring



Health Care



Agriculture



Logistic and Supply Chain



Why large scale wireless network?

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Scalability
Diversity (spatial, temporal)
 Asymptotical Behavior
   Application Requirement
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Presentation Outline



Real World Systems

- 1. OceanSense (2007)
- 2. GreenOrbs (2009-)
- 3. CitySee (2011-)

OceanSense









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Motivation

Silt Deposition problem of Qingdao Port:

- Qingdao port :
 - ➤ one of the ten busiest ports in the world
- Silt Deposition:
 - ≻ Affect the water depth
 - > High uncertainty and high instant uncertainty (tide, wind, etc.)





OceanSense

Monitor the sea!

- The first sea environment monitoring sensor network system in China
- More than 120 sensor nodes
- Temperature, Light, Sea depth



Deployment

• Deployed in the Yellow sea near Qingdao, China



GreenOrbs

http://www.greenorbs.org/















Motivation





Canopy closure estimates





GreenOrbs

✤ Go to the wild!

- Supporting forestry research and applications
- Multiple deployments, each>330 sensor nodes
- Temperature, Light, CO2



Deployment: Overview

Place	Area	Duration	Battery	Scale	Network Diameter	Duty Cycle	Data Volume
University woodland #1	20,000 m ²	1 month (2008)	800 mAh 1.5V	50	6 hops	No	15 Mbytes
University woodland #2	20,000 m ²	10 months (2009)	2200 mAh 1.2V	120	10 hops	5%	272 Mbytes
University woodland #2 and #3	40,000 m ²	1 year (2009.12~)	~8000mAh, 1.5V	330	12 hops	8% or No	140 Mbytes
Tianmu Mountain	200,000 m ²	1.5 months (2009)	~8000mAh, 1.5V	50	10 hops	5%	3 Mbytes
Tianmu Mountain	200,000 m ²	1.5 year (2009.10~)	~8000mAh, 1.5V	200	~ 20 hops	5%	10 Mbytes

TianMu

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Campus



Deployment: Nodes in the Wild









CitySee City-Wide Urban Sensing













Motivation: Global Warming

- Starting from Global Climate Changes
 - Emission of large volume of greenhouse gases is the main reason for global warming
 - ≻CO2, N2O, CH4, HFCs, PFCs, SF6
 - The most greenhouse gases is CO2
 - CO2 generation of human activities: in the city



CitySee

- ✤ Back in the city!
 - Large scale indoor/outdoor environment monitoring
 - More than 1200 sensor nodes
 - Temperature, Light, CO2
 - Mesh routers







System Architecture



Deployment: Locations

Cover more than 1.2 KM² urban area of the Wuxi City



Thermal Power Plant





High emission Factories



Development Zone



Residential Area



Railway Station



Water Source

System Deployment



Deployment: Nodes Deployed

	Normal node	Carbon node	Mesh node	
Microcontroller	MSP430f1611	MSP430f1611	ARM7	
Type of sensor reading	Temperature, humidity, light.	CO ₂	N/A	
Radio module	IEEE 802.15.4 CC2420 2.4GHz	IEEE 802.15.4 CC2420 2.4GHz	IEEE 802.11b NetCard 5.8GHz	
Communication range (m)	150~200	150~200	5000~6000	
Power	2 AA batteries (3V)	12V Rechargeable battery	110V~220V AC	
Power consumption - sleeping (mW)	0.6~1.2	2.4~4.8	N/A	
Power consumption - sensing (mW)	60~90	~2160	N/A	
Power consumption - communication (mW)	60~90	60~90	3000~25000	
Manufactory cost (USD)	~80	~260	~800	



Presentation Outline



LESSONS

System that work in labs fails horribly in practice

- OceanSense:
 - System run out of battery in a week (labs run in months)
 - \succ Nodes destroyed by water
 - > Devices stolen by people: they are interested in the sticks!
- GreenOrbs:
 - ≻ Nodes destroyed by flooding
- CitySee:
 - Installation needs the coordination of various government departments
 - ➢ Require nice encapsulation

- Encapsulation? Encapsulation!
 - Solutions to many of the previously mentioned problems
 - OceanSense:
 - ➤ Waterproof, considering factors such as tide, wind, etc.
 - GreenOrbs:
 - > Waterproof, allow accurate collection of humidity and luminosity
 - CitySee:
 - ➤ Made the nodes nice-looking!









Need good visualization tools

- Allow diagnostics
- Easy to interpret the data and locate the problems



Deployment

- Balance between accuracy, coverage, sustainability and cost
- Regions that doesn't allow deployment
 - > Not allowed by the nature (physically infeasible)
 - ≻ Not allowed due to bad signal (interference, obstruction, etc)
 - ➢ Forbidden by the government
- Deployment challenges: bamboo, large trees





Control the cost!

- Cost reduction is a must when you need so many nodes
 - \succ Node cost
 - \succ Labor cost



NETWORK OBSERVATIONS: GREENORBS

Traffic distribution : balanced in CTP?





The traffic distribution is relatively stable over time

Causes of Packet Losses

- Packet Delivery Ratio (PDR) about 85%
 - Link loss (61%) vs. Node drops (39%)
- Faulty behavior on forwarding nodes



Cumulative distribution of packet loss Causes of packet drops on sensor nodes

Packet Loss Diagnosis



- The green nodes with PRR > 90%.
- The red nodes with PRR < 90%,
- The radius indicates the number of lost packets

Packet Losses: Non-ACK

✤ 84,030 packet loss due to non-ack

- 46.2% of total losses
- 68,444 caused by physical environment (bad links)



Packet Losses: Non-ACK

✤ 84,030 packet loss due to non-ack

- 46.2% of total losses
- 4,361 caused by interferences (contention <--reboot, loop)



Packet Losses: Corrupted Packets

✤ 9,511 corrupted packets

- 9037 real losses (after consider retransmission)
- $\sim 5\%$ of total loss


Packet Losses: Routing Loop

\$5,178 packet loss due to overflow from routing loop - 2.9% of total losses

– 93% of overflow events did not result in packet loss



Packets Loss Summary

Root cause	%
1. sink-side failure	12.5%
1.1 vertical banding	12.45%
2. corruption	5%
3. overflow drops	2.87%
3.1 loop overflow drops	2.85%
3.2 non-loop overflow drops	0.02%
4. no-ack drops	46.2%
4.1 env-no-ack drops	37.6%
4.2 interference-no-ack drops	2.4%
5. reboot (direct impact on loss)	${\sim}0$

About 35% packet losses are unidentified now.

NETWORK OBSERVATIONS: CITYSEE

Where Packets are Lost ?



Traffic Distribution

- Small portion of "critical nodes", verifies the same finding observed from GreenOrbs
- Traffic dynamics exhibits different pattern, e.g. burstness on some nodes



Traffic

volume

Time

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Sink

Network Topology



Nodes closer to the sink have a more stable topology than nodes that are far away.

Summary

Many challenges to make it

- **1. Sustainable** --- energy efficiency and fault diagnosis?
- **2. Robust** --- co-existence?
- **3.Scalable** ---large scale performance?
- **4. Predictable** ---- under varying environment?

Presentation Outline





ZIMO: CROSS-TECHNOLOGY MIMO TO HARMONIZE COEXISTENCE OF ZIGBEE WITH WIFI

Coexistence in ISM Band

ISM band interferences are pervasive and crowded WiFi signal is the primary and first class passenger



Existing Works mainly protect WiFi signal and mitigate cross technology interference [TIMO]
 Some ZigBee signal protection works need modifications or degrade WiFi [Sensys10, Liang]

Experiment Setup

- Two ZigBee Nodes (TX&RX)
- ✤ WiFi APs are in IIT campus
- ZigBee nodes are configured to receive full spectrum interference in full time scale
- Adding controllable AP for tunable interference



Effect of WiFi interference on ZigBee



✓ Short and frequent WiFi data transmission (i.e., flash) play the main role of WiFi interference on ZigBee.

✓ Power-law like distribution indicates the shorter flashes interfere ZigBee signal with exponentially increasing probability, which is a drastic threat for ZigBee signal.

arabite arread for 215000 biginar.

Effect of WiFi interference on ZigBee



(a) Length statistics

(b) Interval statistics

✓ The WiFi interference is distributed across ZigBee symbols, rather than concentrated on particular positions.
 ✓ We need to resort to the signal processing techniques for fundamental solutions.

fundamental solutions

The weed to report to the Sten broads and reason the

ZIMO: Sink Based Design



Cons No.1: ZIMO has more antennas than WiFi AP (N+1) Cons No.2: ZIMO needs at least one preamble is clear Cons No.3: Can work with one ZigBee and multiple WiFi

Cons No.3: Can work with one ZigBee and multiple WiFi

How ZIMO works?



Channel Coefficient Recovery

- Interference in frequency domain
- Sufficient for decoding, insufficient for accurate signal recovery
- CFO Compensation
 - Well done for preamble, insufficient for whole data scale
 - Extremely large with the increasing packet length

Where are opportunities?

Frequency domain is partially overlapped
Time domain also partially overlapped
Power domain shows significant difference



Where are opportunities?



For channel coefficient: Interpolation is simple, effective

interputation is simple, encouve

Where are opportunities?



Linear regression is accurate enough for CFO

Implementation





- Implement using USRP2 N200
 - IEEE STD 802.15.4, 2 MHz Bandwidth
 - OFDM is 20 MHz Bandwidth
- Real trace driven ZIMO decoding
- No carrier sense and MAC timing control

Experimental Results: Macro Benchmark

Recovery ratio



Experimental Results: Macro Benchmark

Throughput



Zigbee Baseline

WiFi Baseline

Interference patterns

Asymptotical Capacity

Two Capacity metrics - channel

*Channel capacity

– achievable single-hop data rate

EK

 $C = \log(1 + SINR)$

Shannon channel theory

Two Capacity metrics - transport

* transport capacity

– end-to-end multi-hop throughput



Impact factors :

- Network Size
- Networking Models
- *** Inference Models**
- * Traffic Models



a meters

Various Models

* deployment models

- arbitrary networks
- random networks



* network scaling models

- dense networks
- extended networks

* Communication (Interference) models

- the protocol model (PrIM)
- Fixed Range Protocol Model (fPrIM)
- physical model (PhIM)
- generalized physical model (GphM, also called GCM)
- Others

* Traffic models

- Unicast
- Broadcast
- Multicast
- Anycast
- Many-to-one







Results Summary

Milestone Results : Unicast, PrIM



Milestone Results : Unicast, PhIM



Milestone Results : Broadcast, PrIM



Milestone Results : Multicast, PrIM



Our Results : Multicast, PrIM



Li et al, MobiCom 2007, REN, $n_s = \Theta(n)$

Brief Summary

The aggregate multicast capacity of n sessions is

$$\Lambda_{n_d}(n) = \begin{cases} \Theta(\sqrt{\frac{n}{\log n}} \bullet \frac{W}{\sqrt{n_d}}) & \text{when} \quad n_d = O(\frac{n}{\log n}) \\ \Theta(W) & \text{when} \quad n_d = \Omega(\frac{n}{\log n}) \end{cases}$$

- Our results unify previous results
 - **1.** Unicast (when $n_d=2$): $\Theta(\sqrt{\frac{n}{\log n}} \bullet W)$ by Gupta and Kumar
 - **2. Broadcast** (when $n_d = n$): $\Theta(W)$ by Keshavarz-Haddad et al., Mobicom'06
 - **3.** Multicast $(n_s = n^{\epsilon} \text{ and } n_d = n^{1-\epsilon})$, $O(\sqrt{\frac{n}{n_d \log n}})$ by Shakkottai et al., Mobihoc'07

Our Results : Multicast, GCM



Multicast Capacity for **REN**, $n_s = \Theta(n)$, Li et al. [MobiCom 2008]. *Wang*, Li et al. [INFOCOM 2010, 2011].

Protocol Interference Model General Approaches
Multicast under Protocol Model

* Data Copies Argument (upper bound)

- Estimate the expected (or asymptotic lower bound) number of nodes N(b) that received (or listened) a bit b.
- Capacity at most $n \cdot W/N(b)$

 \succ since all nodes receive at rate at most $n \cdot W$.



Upper-bound Proof Flow



Lower-Bound: Routing and Scheduling

✤ Build EMST

- Routing structure using EMST as backbone
- Need to bound the conflict and total data copies
 - ≻ The lower-bound of multicast tree length w.h.p.? EMST?
 - > Maximum number conflicting flows in the network w.h.p
 - Using VC dimension (proved to be $O(\log n_d)$), and VC theorem





Lower-bound Proof Flow



Gaussian Channel Model General Approaches

Multicast under Gaussian Model

- Two kind of links
 - Inside Links
 - Outside Links



Relationship between links

- $\bigstar l_c$: max link length in giant component.
- $\bigstar l_c$: the max distance between any node not in GC and the giant cluster



Upper-bound Proof Techniques 1

★ There is a link *uv*, that will be used by many flows (say *f*) ⇒ the minimum data rate $-\min \lambda_i \leq \text{rate supported by } uv / f$



Upper-bound Proof Techniques 2

- There is an isolated cluster *C* of nodes, and *f* flows will have links going inside this cluster
 - − min λ_i ≤ total rate supported by links reaching *C*/*f*



Lower Bounds Techniques





Low Bound: Routing, Scheduling

- ✤ First build EMST of receivers
- Build highway using cell size 1
 Each highway link data rate O(1)
 - Each highway link data rate O(1)
- ✤ Build second-class highway using cell size (log n)^{1/2}
- Node sends its data to highway (solid lines) by multi-hop second class highway (dashed line)



Our New Techniques

- Parallel Arterial Road Systems
 - longer links to connect isolated nodes to highway



*****Parallel Scheduling

Other Research

Cyber Physical Systems



Cognitive Radio Networks



Our iGaze Glasses



TAGORAM

Drawing in the Air



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MS Students



Students in IIT



Domestic Students



Deployment Videos





Thank you !

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